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Marcus External Contextual Grammars with Choice of the Languages of Primitive and Generalized Primitive Words. An Alternative Proof

To the honor Professor Masami Ito on his 70-th birthday

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Abstract

In this paper we unify some well-known results describing an alternative proof that all of the languages of primitive, quasi-primitive, and hyper-primitive words are Marcus external contextual languages with choice.

Keywords: Formal languages, Marcus contextual languages, combinatorics of words and languages.

1 Preliminaries

Marcus contextual grammars were introduced and intensively studied by S. Marcus and his students (see [11, 13]). The word is primitive if it is not a

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power of its proper prefix. The quasi-primitivity and hyper-primitivity are natural extensions of this concept. The relation of the language of primitive words to the Marcus contextual languages was studied first in [4]. On the line of this research, the relation of the language of quasi-primitive and hyper-primitive words and their certain further generalizations was described in [2] and [6]. In this paper we unify some of the results in [4, 2, 6] describing an alternative proof that all of the languages of primitive, quasi-primitive, and hyper-primitive words are Marcus external contextual languages with choice.

All notion and notations not defined here we refer to [3]. A *word* (over Σ) is a finite sequence of elements of some finite non-empty set Σ. We call the set Σ an *alphabet,* the elements of Σ *letters.* If u and v are words over an alphabet Σ, then their *catenation* uv is also a word over Σ. Especially, for every word u over Σ, $u\lambda = \lambda u = u$, where λ denotes the *empty word.* Given a word u, we define $u^0 = \lambda$, $u^n = u^{n-1}u$, $n > 0$, $u^* = \{u^n : n \geq 0\}$ and $u^+ = u^* \setminus \{\lambda\}$.

For every triplet u, v, w of words we say that u is a *prefix*, w is a *suffix*, and v is a *subword* of uvw. If u (v, w) is nonempty then we speak about *proper prefix* (*proper subword*, *proper suffix*). A word z is called *overlapping* or *bordered* if there are $u, v, w \in \Sigma^+$ with $z = uw = wv$.

The *length* $|w|$ of a word w is the number of letters in w, where each letter is counted as many times as it occurs. Thus $|\lambda| = 0$. By the *free monoid* Σ^* *generated by* Σ we mean the set of all words (including the *empty word* λ) having catenation as multiplication. We set $\Sigma^+ = \Sigma^* \setminus \{\lambda\}$, where the subsemigroup Σ^+ of Σ^* is said to be the *free semigroup generated by* Σ. Subsets of Σ^* are referred to as *languages* over Σ.

A *primitive word* (over Σ, or actually over an arbitrary alphabet) is a nonempty word not of the form w^m for any nonempty word w and integer $m \geq 2$. The set of all primitive words over Σ will be denoted by $Q(\Sigma)$, or simply by Q if Σ is understood. Q has received special interest: Q and $\Sigma^+ \setminus Q$ play an important role in the algebraic theory of codes and formal languages (see [7, 8, 9, 14]). If $u \in \Sigma^+$ can not be written into the form $u = v^n v', n \geq 2$ such that $u, v \in \Sigma^+$ and v' is a prefix of u then we say that u is *strongly-primitive.*

We say that a word $u \in \Sigma^+$ *is covered* by the word $v \in \Sigma^+$ if for every $u', u'' \in \Sigma^*, a \in \Sigma$ with $u = u'au''$ there are $v_1, v_2, v_3, v_4 \in \Sigma^*$ with $u = v_1v_2av_3v_4, v = v_2av_3, u' = v_1v_2, u'' = v_3v_4$.

A word $u \in \Sigma^+$ is called *hyper-primitive* if it can not be covered by any of its proper subwords.

$u \in \Sigma^+$ is *super strongly primitive* if $u \neq v^n v', n \geq 2$ such that v has a suffix v'' for which $v''v'$ is a prefix of u.

u is called *strongly hyper-primitive* if $u \neq wv'$, where w is covered by v, which is one of its proper subwords, and v' is a prefix of v.

Finally, u is *hyper hyper-primitive* if $u \neq wv'$, where w is covered by v, which is one of its proper subwords, and w has a suffix v'' such that $v''v'$ is a prefix of v.

Denote, in order, $SQ(\Sigma), HQ(\Sigma), SSQ(\Sigma), SHQ(\Sigma), HHQ(\Sigma)$, or, if Σ is understood, then SQ, HQ, SSQ, SHQ, HHQ the language of all strongly primitive, hyper primitive, super strongly primitive, strongly hyper-primitive, and hyper hyper primitive words (over Σ).

Moreover, denote by $|H|$ the *cardinality* of H for every set H.

A *(Marcus) contextual grammar with choice* is a structure $G = (V, A, C, \varphi)$, where V is an alphabet, A is a finite language over V, C is a finite subset of $V^* \times V^*$, and $\varphi : V^* \to 2^C$. If $\varphi(x) = C$ holds for every $x \in V^*$ then we say that G is a (Marcus) contextual grammar without choice and then we omit φ sometimes.

We define two relations on V^* as usual: for any $x \in V^*$, we write

$x \Rightarrow_{ex} y$ if and only if $y = uxv$, for a context (u, v) in $\varphi(x)$,
$x \Rightarrow_{int} y$ if and only if $x = x_1x_2x_3, y = x_1ux_2vx_3$ for any $(u, v) \in \varphi(x_2)$.

Denote $\overset{*}{\Rightarrow}_{ex}$, $\overset{*}{\Rightarrow}_{in}$ the reflexive and transitive closure of these relations and let $L_\alpha(G) = \{x \in V^* : w \overset{*}{\Rightarrow}_\alpha x, w \in A\}$ for $\alpha \in \{ex, in\}$. Then $L_{ex}(G)$ is the *(Marcus) external contextual language (with or without choice) generated by* G, and similarly, $L_{in}(G)$ is the *(Marcus) internal contextual language (with or without choice) generated by* G. Now let $G = (V, A, \varphi)$, where V is an alphabet, A is a finite language over V, C is a finite subset of $V^* \times V^*$, and $\varphi : V^* \times V^* \times V^* \to 2^C$.[1]

Define the relation $\Rightarrow$ on V^* such that $x \Rightarrow y$ for some $x, y \in V^*$ if and only if $x = x_1x_2x_3, y = x_1ux_2vx_3, x_1, x_2, x_3 \in V^*, (u, v) \in \varphi(x_1, x_2, x_3)$. Moreover, let $\overset{*}{\Rightarrow}$ denote the reflexive and transitive closure of $\Rightarrow$. Thus $L(G)$ is defined to be a *(Marcus) total contextual grammar (with or without choice) generated by* G. If $\varphi(x_1, x_2, x_3) = C$ holds for every $x_1, x_2, x_3 \in V^*$ then we say that G is a (Marcus) total contextual grammar without choice and sometimes we omit φ having this property.

[1]Observe that the definition of φ is not the same as before.

The following statement is a unified form of some results in [2, 4, 6]. It has been formulated by [6].

Theorem 1 *[2, 4, 6] The languages Q, SQ, and HQ are external contextual languages with choice. This is not true for the sets SSQ, SHQ, and HHQ, furthermore, none of the sets Q, SQ, HQ, SSQ, SHQ, and HHQ is an external contextual language without choice or an internal contextual language with or without choice.* □ □

We shall use the following results.

Theorem 2 *[5] Let $u, v \in \Sigma^+, s, t \geq 1$, with $s \neq t$. If $\sqrt{u} \neq \sqrt{v}$ and $uv^s \notin Q$, then $uv^t \in Q$.* □

Theorem 3 *[1] Let $u, v \in Q, u^m = v^k w, k, m \geq 2$ for some prefix w of v. Then $u = v$ and $w \in \{u, \lambda\}$.* □ [2]

Theorem 4 *[14][BorweinLemma] Let $u \in \Sigma^+, u \notin a^+, a \in \Sigma$. Then at least one of ua, u must be primitive.* □

Theorem 5 *[10] If $uv = vq, u \in \Sigma^+, v, q \in \Sigma^*$, then $u = wz, v = (wz)^k w, q = zw$ for some $w \in \Sigma^*, z \in \Sigma^+$ and $k \geq 0$.* □

We shall use the following two widely known consequences of Theorem 5.

Proposition 6 *For every bordered word $z \in \Sigma^+$ there exists a nonempty word $u \in \Sigma^+$ and a (not necessarily nonempty) word $v \in \Sigma^*$ having $z = uvu$.* □

Theorem 7 *[10] Let $u, v \in \Sigma^+$ with $uv = vu$. There exists $w \in \Sigma^+$ with $u, v \in w^+$.* □

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[2] This statement can also be derived directly from [5].

2 Results

Next we show alternative proofs of some known results.

Theorem 8 *[2, 6] Let V be an alphabet with $|V| \geq 2$. If $awb \in SQ$ where $w \in V^*$ and $a, b \in V$, then $aw \in SQ$ or $wb \in SQ$.*[3]

Proof: Suppose the contrary. Then $aw, wb \in SPer$, i.e., there are $u, v \in V^+, u', v' \in V^*$, positive integers $m, n \geq 2$ such that u' is a prefix of u, v' is a prefix of v, and $u^m u' = aw, v^n v' = wb$.

Then $u = aw_1w_2$ and $u' \in \{\lambda, aw_1\}$ for some $w_1, w_2 \in V^*$. Similarly, $v = w_3bw_4$ and $v' \in \{\lambda, w_3b\}$ for an appropriate pair $w_3, w_4 \in V^*$. Thus we can write $w \in \{(w_1w_2a)^m w_1, (w_1w_2a)^{m-1}w_1w_2, (w_3bw_4)^n w_3, (w_3w_4b)^{n-1}w_3w_4\}$. Let, say, $w = (w_1w_2a)^m w_1 = (w_3bw_4)^n w_3$. By the symmetricity we may assume $|w_1| \leq |w_3|$. Thus $(w_1w_2a)^m = (w_3bw_4)^n w'$ for some prefix w' of w_3. Applying Theorem 3, $\sqrt{w_1w_2a} = \sqrt{w_3bw_4}$. Therefore, $w_3bw_4 = w_3w''a$ for some $w'' \in V^*$. Hence $awb = a(w_3bw''a)^n w_3b = (aw_3bw'')^n aw_3b \notin SQ$, a contradiction.

We can get the same conclusion if $w = (w_1w_2a)^m w_1 = (w_3w_4b)^{n-1}w_3w_4$ and $n > 2$ (or $w = (w_1w_2a)^{m-1}w_1w_2 = (w_3bw_4)^n w_3$ and $m > 2$). Thus let $w = (w_1w_2a)^m w_1 = w_3w_4bw_3w_4$ (with $n = 2$). But then $(w_1w_2a)^m w_1 b = (w_3w_4b)^2$ with $m \geq 2$. Applying again Theorem 3, $\sqrt{w_1w_2a} = \sqrt{w_3w_4b}$ with $a = b$. Therefore, $awb = aw_3w_4bw_3w_4b = (aw_3w_4)^2 a \notin SQ$, a contradiction.

We can derive the impossibility of $w = w_1w_2aw_1w_2 = (w_3bw_4)^n w_3$ and $n \geq 2$ in the same way.

The rest of the cases is the equality $w_1w_2aw_1w_2 = w_3w_4bw_3w_4$. But then $|w_1w_2| = |w_3w_4|$ which implies $w_1w_2a = w_3w_4b$, i.e., $a = b$. Then $awb = aw_3w_4bw_3w_4b = (aw_3w_4)^2 a \notin SQ$, a contradiction again. □

Lemma 9 *If a word w can be covered by a word va, with $v \in \Sigma^*, a \in \Sigma$, then vb is not a subword of w, for any $b \in \Sigma, b \neq a$.*

Proof: Consider a covering of w by va. We will assume that vb can occur in w and show that it leads to a contradiction.
There are two possibilities for vb to occur in w:

Case 1. vb is a proper subword (not only pre- or suffix) of $v'v$, where v' is a prefix of v: in this case vb is neither a prefix nor a suffix of $v'v$ because

[3] $a = b$ is possible.

$va \neq vb$. Thus v has two different borders, i.e. by Proposition 6, $v = x_1ux_1$ and $v = x_2y'x_2$. Without loss of generality we can assume $|x_2| < |x_1|$. Then x_1 itself is bordered, hence, applying Proposition 6 again, $x_1 = xyx$, for some x, y. This gives us $v = xyxuxyx$ and because v overlaps twice with itself (by xyx and also by x), $v = xyxuxyx = xuxyxz$, for some z, but then x is a suffix of z and immediately before it is y, so $xyxuxyx = xuxyxyx$. Simplifying gives us $yxu = uxy$, hence

$$xuxyxyx = xyxuxyx = xyxyxux \text{ with } v = xyxuxyx, \tag{1}$$

taking away the first x, we get $uxyxyx = yxyxux$, so $ux = (yx)^2$. Therefore, by Theorem 7, $ux, (yx)^2 \in z^+$ for some $z \in \Sigma^+$. From here applying (1), $v = xz^k$, where z is a primitive word and $k \geq 3$. Moreover, since x is a suffix of v, we get $x = z'z^j$, with z' a suffix of z and $j < k$, so $z = z''z'$ and $v = z'(z''z')^{j+k}$, with $z''z'$ primitive, therefore $z'(z''z')^{j+k}b$ would have to be a proper subword of either $z'(z''z')^{j+k}az'(z''z')^{j+k}$ or $z'(z''z')^i$, with $i > j+k$. In both cases the first letter of z'' would have to be at the same time a and b, contradiction.

Case 2. vb is a proper subword of vav. In this case v from va overlaps the first v in vav with a part u_1 and the second with u_2, that is, $v = u_1au_2$ and $v = u_2bu_1$. If $|u_1| = |u_2|$, we instantly get $a = b$, contradiction. Without loss of generality $|u_1| < |u_2|$, and then u_1 is a border of u_2 so, applying Theorem 5, for some $x \in \Sigma^*, y \in \Sigma^+$ we have $u_1(= (xy)^ix) = x(yx)^i$ and $u_2(= (xy)^jx) = x(yx)^j$, with $1 \leq i < j$. This gives $v = x(yx)^iax(yx)^j = x(yx)^jbx(yx)^i$. Taking away $x(yx)^i$ from both sides we get $ax(yx)^{j-i} = x(yx)^{j-i}b$. By this equality, $x \neq \lambda$ implies $ax = xc$ and $dx = xb$ for some $c, d \in \Sigma$. Hence we could get $x \in a^+ \cap b^+$, a contradiction. Therefore, $x = \lambda$. Then $ay^{j-i} = y^{j-i}b$ with $a \neq b$ and $i < j$. (By $a \neq b$, $i = j$ would be impossible even if we would not suppose before $i < j$.) By this connection, $y \neq \lambda$ implies $ay = yc$ and $dy = yb$ for some $c, d \in \Sigma$. Then $y \in a^+ \cap b^+$, which is impossible unless $a = b$. □

Theorem 10 *[6] For any word w and (not necessarily distinct) letters $a, b \in \Sigma$, if $aw, wb \notin HQ$, then $awb \notin HQ$.*

Proof: If $aw \notin HQ$, then there is some hyper-primitive av which covers aw. Similarly, there is some hyper-primitive ub which covers wb. Without loss of generality, we can assume $|v| \leq |u|$. Then, u is a suffix of v, therefore wherever there is an occurrence of v in the string, it ends in u. Now, Lemma 9 tells us that if ub covers wb, and $c \neq b$, then uc is not a subword of wb.

There are two cases.

Case 1. $a \neq b$. Whenever v appears in the string wb, it should be followed by b. From here, we get that avbcovers awb, so $awb \notin HQ$.

Case 2. $a = b$. Whenever v appears in the string wa, it should be followed By a. From here, we get that ava covert awa, so $awa \notin HQ$. □

Corollary 11 *Let V be an alphabet with $|V| \geq 2$. If $awbc \in XQ$, where $XQ \in \{Q, SQ, HQ\}$, $w \in V^*$ and $a, b, c \in V$[4], then one of aw, awb, wbc is in XQ.*

Proof: If $XQ = Q$ and $awb \notin a^+$, then Theorem 4 implies that one of aw, awb should be in Q. If $XQ = Q$ and $awb \in a^+$, then $awbc \in XQ$ implies $c \neq a$. In this case, $wbc \in a^+c$ with $a \neq c$, for which $wbc \in Q$ obviously holds. If $XQ \in \{SQ\}$ then by Theorem 8, if $XQ \in \{SQ, HQ\}$ then by Theorem 10 we have that one of awb, wbc should be in XQ. □

On the basis of Lemma 11, similarly to Theorem 12 published by [6], next we show an alternative (and unified) proof of the next statement which is a union of three previous results.

Theorem 12 *[2, 4, 6] All of the languages Q, SQ, HQ are Marcus external contextual languages with choice.*

Proof: Let $G = (V, A, C, \varphi)$ be be an external Marcus contextual grammar with choice defined by $A = V$, $C = \{(\lambda, \lambda), (\lambda, a), (\lambda, ab), (a, \lambda) : a, b \in V\}$, moreover, let for every $w \in V^*$, $z \in \varphi(w)$ with

$$z = \begin{cases} \{(\lambda, \lambda)\} & \text{if } |V| = 1, \\ \{(a, \lambda)\} & \text{if } a \in V \text{ and } aw \in XQ, \\ \{(\lambda, a)\} & \text{if } a \in V \text{ and } wa \in XQ, \\ \{(\lambda, ab)\} & \text{if } a, b \in V \text{ and } wab \in XQ.^5 \end{cases}$$

Moreover, let $XQ \in \{Q, SQ, HQ\}$. Obviously, the proposition holds true for $|V| = 1$. Hence we assume $|V| \geq 2$. By the definition of the grammar G, it is obvious that $L_{ex}(G) \subseteq SQ$. Now we prove that $SQ \subseteq L_{ex}(G)$

[4] a, b, c are not necessarily distinct

by induction. By definition, $V \cap SQ = V(= A)$ and $V^2 \cap SQ = \{ab \mid a, b \in V, a \neq b\}$. Similarly, $V \cap V^3 = \{abc \mid a, b, c \in V, a \neq b, a \neq c, b \neq c\} \cup \{aab, abb \mid a, b \in V, a \neq b\}$. Moreover, by our construction, $a, b \in V$ and $a \neq b$ imply $a \Rightarrow_{ex} ab$. Thus we have $(V \cup V^2) \cap Q \subseteq L_{ex}(G)$. Similarly, by our construction, $a, b, c \in V$ and $a \neq b, a \neq c, b \neq c$ imply $ab \Rightarrow_{ex} abc$ and $a, b \in V$ and $a \neq b$ imply $ab \Rightarrow_{ex} abb$ and $ab \Rightarrow_{ex} aab$. Now, assume that $(V \cup V^2 \cup \cdots \cup V^n) \cap XQ \subseteq L_{ex}(G)$ for some $n \geq 3$. Let $u \in V^{n+1} \cap XQ$ and let $u = awbc \in XQ$ where $a, b, c \in V$. (Note that a, b, c are not necessarily distinct.) Corollary 11 states that, by this condition, one of aw, awb, wbc in XQ. Hence, either $aw \in XQ$ with $aw \Rightarrow_{ex} awbc$ or $awb \in XQ$ with $awb \Rightarrow_{ex} awbc$, or $wbc \in XQ$ with $wbc \Rightarrow_{ex} awbc$. □

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